

Operating System

Lecture 4 /

Chapter 4 (Threads)

Multicore Programming

Multithreading Models

Thread Libraries

Implicit Threading

Threading Issues



Process Concept
Process Scheduling
Operations on Processes
Interprocess Communication



Threads

A thread is a basic unit of CPU utilization; it comprises a thread ID, a program counter, a register set, and a stack.

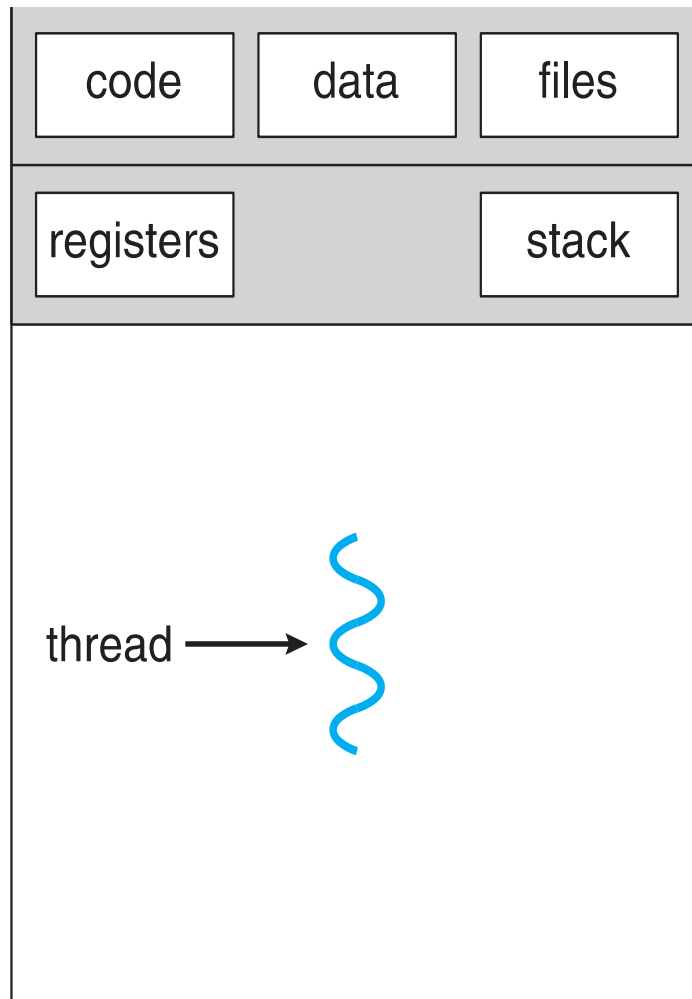
It **shares** with other threads belonging to the same process its code section, data section, and other operating-system resources, such as open files and signals.

A traditional (or **heavyweight**) process has a single thread of control.

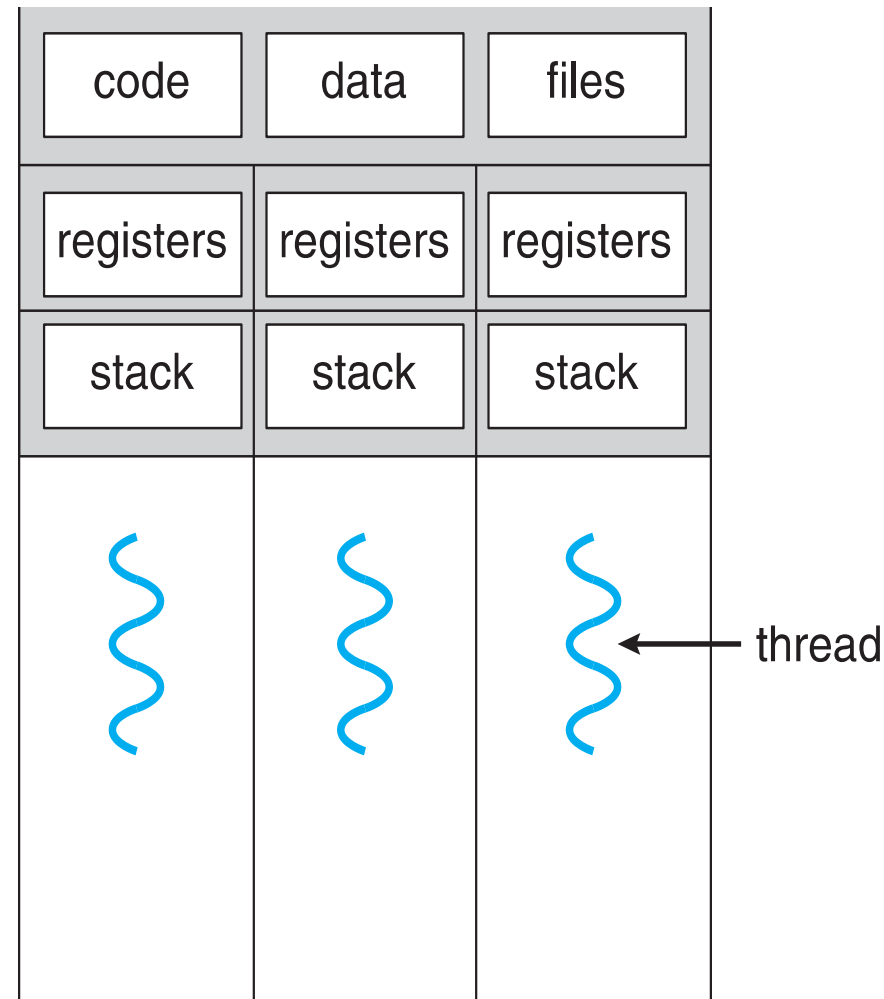
If a process has multiple threads of control, it can perform more than one task at a time.



Single and Multithreaded Processes



single-threaded process



multithreaded process

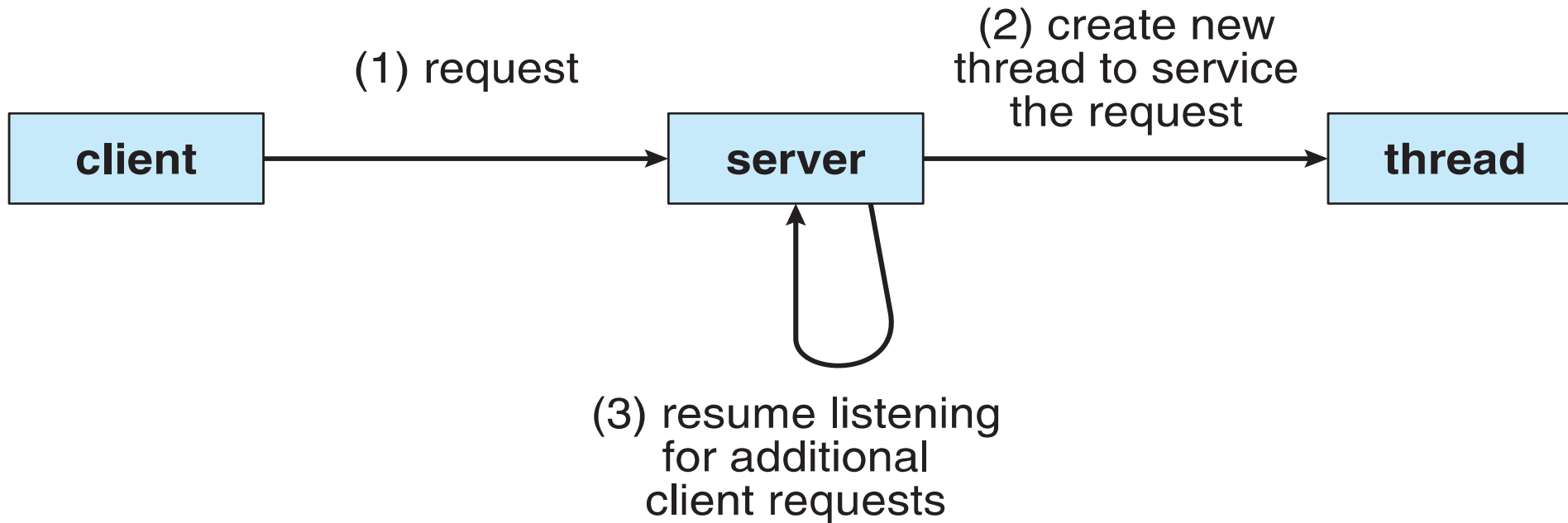


Threads Motivation

- Most modern applications are multithreaded
- Threads run within application
- Multiple tasks with the application can be implemented by separate threads
 - Update display
 - Fetch data
 - Spell checking
 - Answer a network request
- Process creation is heavy-weight while thread creation is light-weight
- Can simplify code, increase efficiency
- Kernels are generally multithreaded



Multithreaded Server Architecture



If the web-server process is multithreaded, the server will create a separate thread that listens for client requests.

When a request is made, rather than creating another process, the server creates a new thread to service the request and resume listening for additional requests.

Threads Benefits

- **Responsiveness** – may allow continued execution if part of process is blocked, especially important for user interfaces
- **Resource Sharing** – threads share resources of process, easier than shared memory or message passing
- **Economy** – cheaper than process creation, thread switching lower overhead than context switching
- **Scalability** – process can take advantage of multiprocessor architectures



Multicore Programming

Multicore or **multiprocessor** systems putting pressure on programmers, challenges include:

- Dividing activities**

- Balance**

- Data splitting**

- Data dependency**

- Testing and debugging**

Parallelism implies a system can perform more than one task simultaneously

Concurrency supports more than one task making progress

- Single processor / core, scheduler providing concurrency

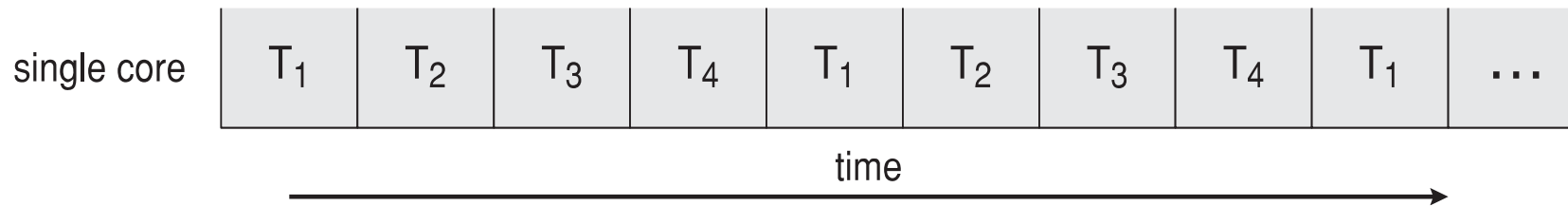


- Types of parallelism
 - **Data parallelism** – distributes subsets of the same data across multiple cores, same operation on each
 - **Task parallelism** – distributing threads across cores, each thread performing unique operation
- As # of threads grows, so does architectural support for threading
 - CPUs have cores as well as *hardware threads*
 - Consider Oracle SPARC T4 with 8 cores, and 8 hardware threads per core

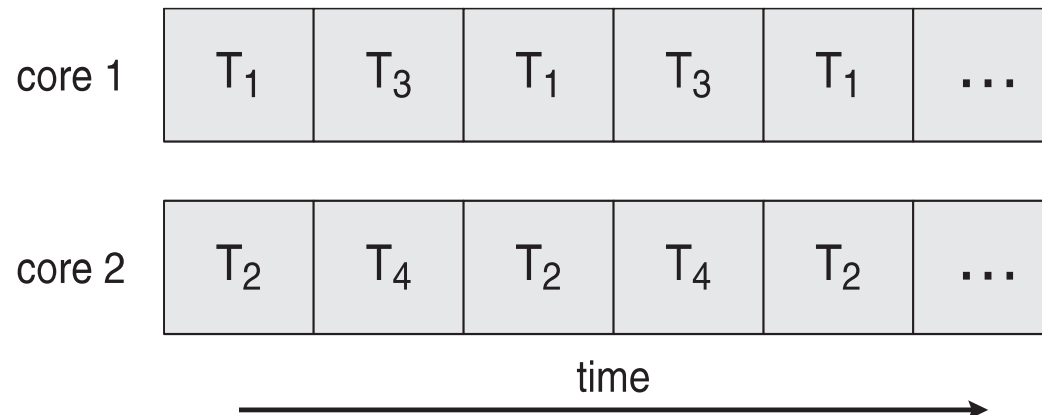


Concurrency vs. Parallelism

■ Concurrent execution on single-core system:



■ Parallelism on a multi-core system:



User Threads and Kernel Threads

- **User threads** - management done by user-level threads library
- Three primary thread libraries:
 - POSIX **Pthreads**
 - Windows threads
 - Java threads
- **Kernel threads** - Supported by the Kernel
- Examples – virtually all general purpose operating systems, including:
 - Windows
 - Solaris
 - Linux
 - Tru64 UNIX
 - Mac OS X

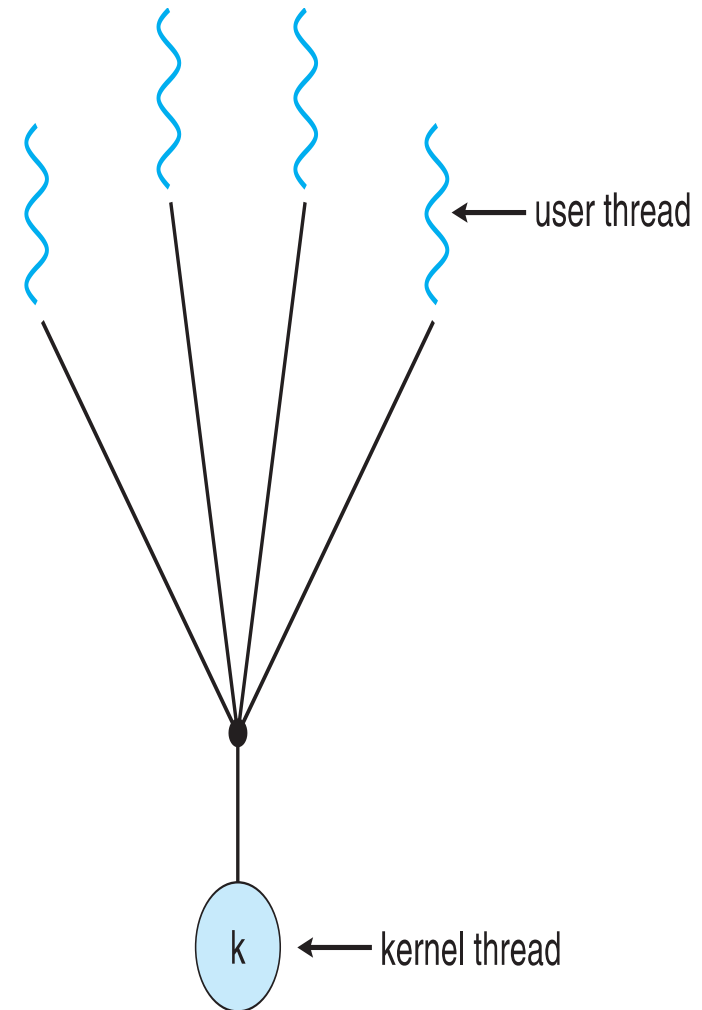


- Many-to-One
- One-to-One
- Many-to-Many

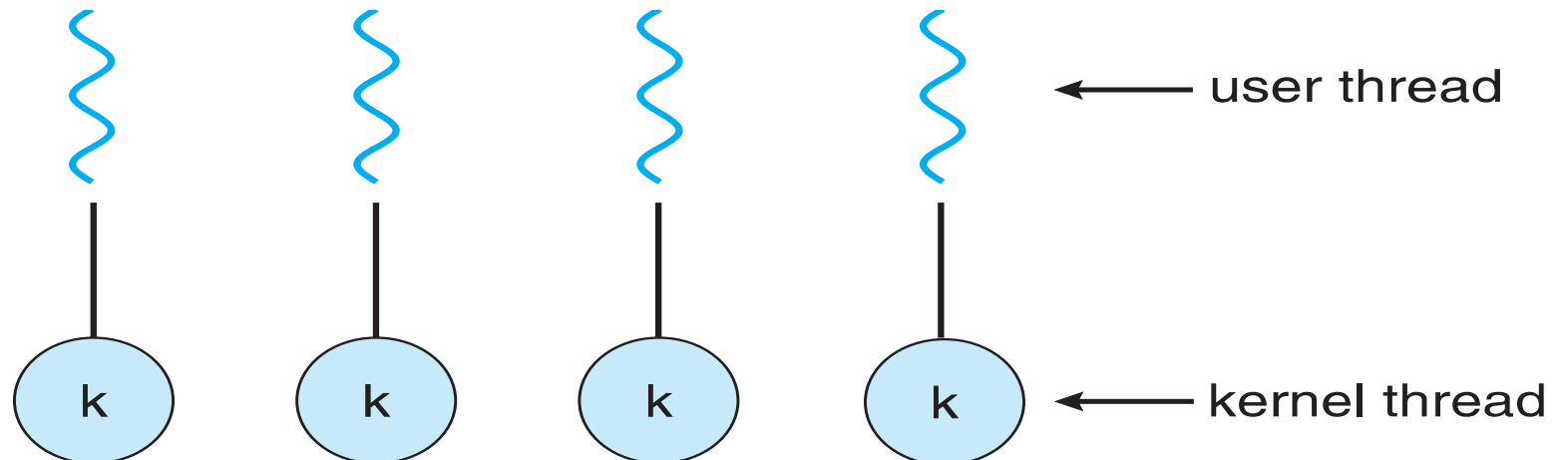


Many-to-One

- Many user-level threads mapped to single kernel thread
- One thread blocking causes all to block
- Multiple threads may not run in parallel on multicore system because only one may be in kernel at a time
- Few systems currently use this model
- Examples:
 - **Solaris Green Threads**
 - **GNU Portable Threads**

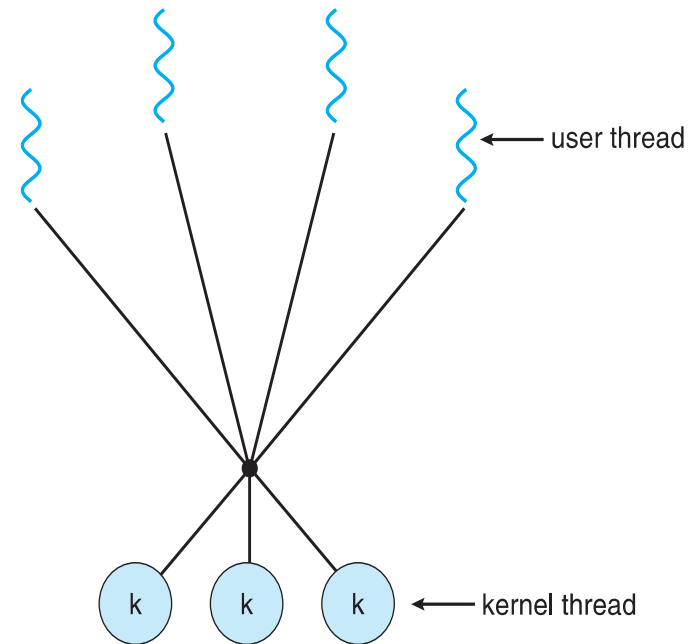


- Each user-level thread maps to kernel thread
- Creating a user-level thread creates a kernel thread
- More concurrency than many-to-one
- Number of threads per process sometimes restricted due to overhead
- Examples
 - Windows
 - Linux
 - Solaris 9 and later

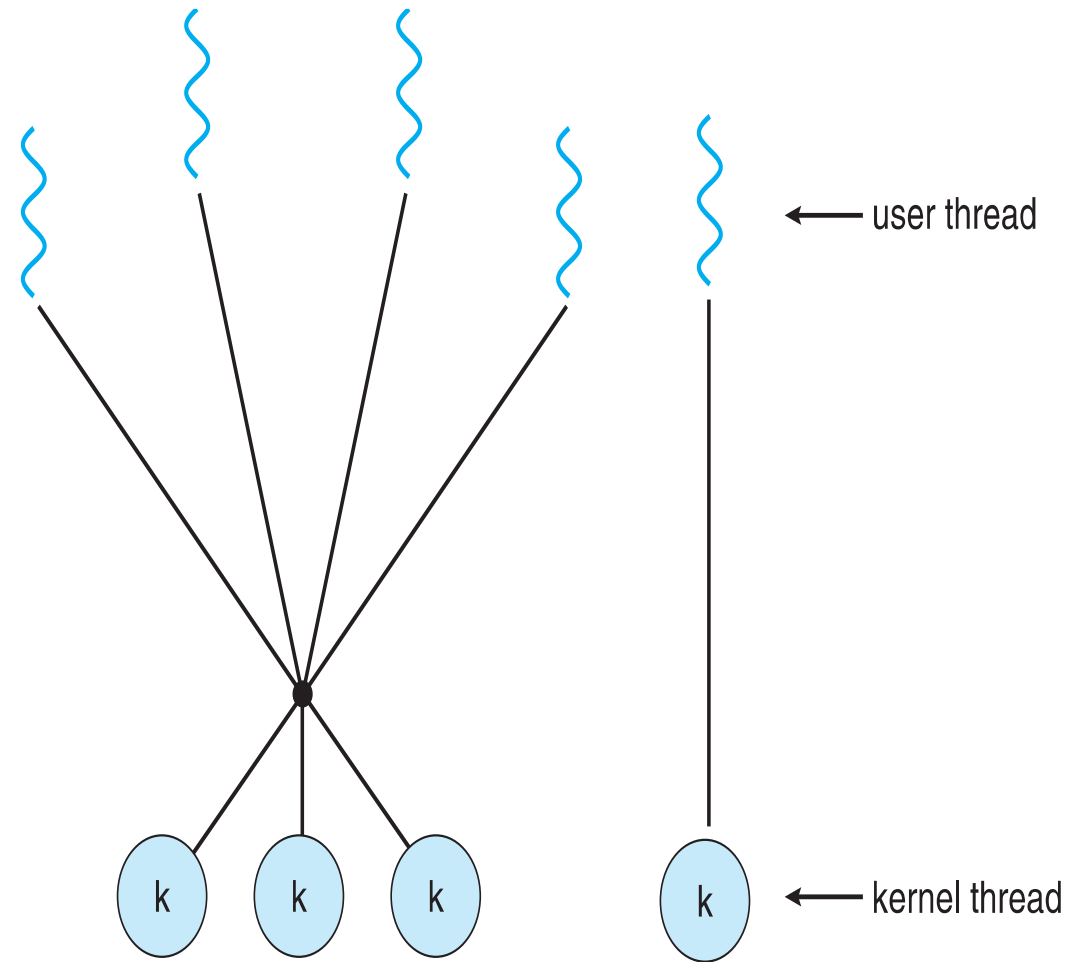


Many-to-Many Model

- Allows many user level threads to be mapped to many kernel threads
- Allows the operating system to create a sufficient number of kernel threads
- Solaris prior to version 9
- Windows with the *ThreadFiber* package



- Similar to M:M, except that it allows a user thread to be **bound** to kernel thread
- Examples
 - IRIX
 - HP-UX
 - Tru64 UNIX
 - Solaris 8 and earlier



- **Thread library** provides programmer with API for creating and managing threads
- Two primary ways of implementing
 - Library entirely in user space
 - Kernel-level library supported by the OS



- General strategies for creating Multiple Threads:
 - Asynchronous
 - Synchronous
- Synchronous Threading Libraries
 - Pthreads
 - Windows Threads
 - Java Threads



- Growing in popularity as numbers of threads increase, program correctness more difficult with explicit threads
- Creation and management of threads done by compilers and run-time libraries rather than programmers
- Three methods explored
 - Thread Pools
 - OpenMP
 - Grand Central Dispatch
- Other methods include Microsoft Threading Building Blocks (TBB), **java.util.concurrent** package



- Create a number of threads in a pool where they await work
- Advantages:
 - Usually slightly faster to service a request with an existing thread than create a new thread
 - Allows the number of threads in the application(s) to be bound to the size of the pool
 - Separating task to be performed from mechanics of creating task allows different strategies for running task
 - i.e.Tasks could be scheduled to run periodically
- Windows API supports thread pools:



- Semantics of **fork()** and **exec()** system calls
- Signal handling
 - Synchronous and asynchronous
- Thread cancellation of target thread
 - Asynchronous or deferred
- Thread-local storage
- Scheduler Activations



OS Kernel Assignment



Submission dead line 2018 23:55

Next Thursday No exceptions would be accepted

Chapter 6

CPU scheduling

